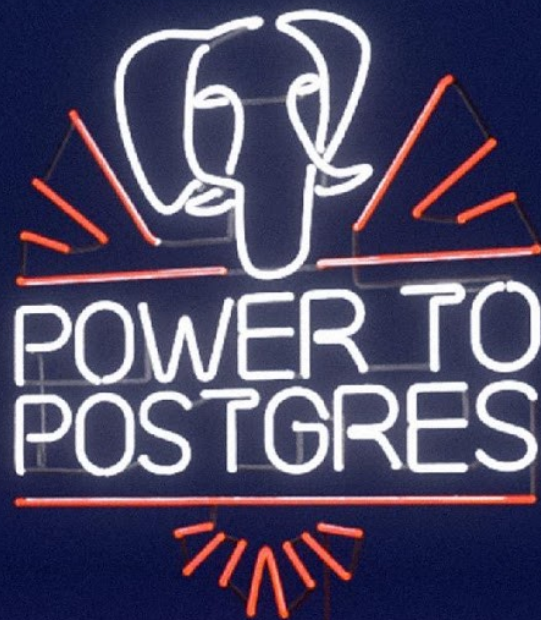


# Don't Do This

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# What is this talk?

- Not all-inclusive
- There is literally nothing you cannot mess up
- Misconceptions
- Confusing things
- Common but impactful mistakes



# We'll be looking at

- Bad SQL
- Improper data types
- Improper feature usage
- Performance considerations
- Security considerations



# Bad SQL

# NOT IN

(i)

Doesn't work the way you expect!

- As in: `SELECT ... WHERE ... NOT IN (SELECT ...)`
- SQL is not Python or Ruby!
  - `SELECT a FROM tab1 WHERE a NOT IN (1, null);` returns NO rows!
  - `SELECT a FROM tab1 WHERE a NOT IN (SELECT b FROM tab2);` same, if any b is `NULL`
- Why is this bad even if no `NULLs`?
  - Query planning / optimization
  - Subplan instead of anti-join



# NOT IN

## What to do instead?

- Anti-join
- ```
SELECT col
FROM tab1
WHERE NOT EXISTS
  (SELECT col
   FROM tab2
   WHERE tab1.col = tab2.col);
```

(ii)



# NOT IN

(iii)

Or:

- ```
SELECT col
FROM tab1
LEFT JOIN tab2 USING (col)
WHERE tab2.col IS NULL;
```
- **NOT IN** is OK, if you know there are no **NULLs**
  - e.g. excluding constants: **NOT IN (1,3,5,7,11)**



# BETWEEN

(i)

Especially with **TIMESTAMPS**

- **BETWEEN (1 AND 100)** is inclusive (closed interval)
- When is this bad?

```
SELECT sum(amount)
FROM transactions
WHERE transaction_timestamp
BETWEEN ('2023-02-05 00:00' AND '2023-02-06 00:00');
```





# BETWEEN

(ii)

Be explicit instead, and use:

```
SELECT sum(amount)
FROM transactions
WHERE transaction_timestamp >= '2023-02-05 00:00'
AND transaction_timestamp < '2023-02-06 00:00';
```



# Using upper case in identifiers

## For table or column names

- Postgres makes everything lower case unless you double quote it
- `CREATE TABLE Plerp (...);`  
`CREATE TABLE "Quux" (...);`
  - Creates a table named `plerp` and one named `Quux`
  - `TABLE Plerp;` works – `TABLE "Plerp";` fails
  - `TABLE Quux;` fails – `TABLE "Quux";` works
  - Same with column names
- For pretty column names: `SELECT col FROM plerp AS "Pretty Name";`



# Improper data types

# TIMESTAMP (WITHOUT TIME ZONE)

## a.k.a. naïve timestamps

- Stores a date and time with no time zone information
  - Arithmetic between timestamps entered at different time zones is meaningless and gives wrong results
- **TIMESTAMPTZ (TIMESTAMP WITH TIME ZONE)** stores a moment in time
  - Arithmetic works correctly
  - Displays in your time zone, but can display it **AT TIME ZONE**
- Don't use **TIMESTAMP** to store UTC because the DB doesn't know it's UTC



# TIMETZ

Or **TIME WITH TIME ZONE** has questionable usefulness

- Only there for SQL compliance
  - Time zones in the real world have little meaning without dates
  - Offset can vary with Daylight Savings
  - Not possible to do arithmetic across DST boundaries
- Use **TIMESTAMPTZ** instead



# CURRENT\_TIME

Is `TIMETZ`. Instead use:

- `CURRENT_TIMESTAMP` or `now()` for a `TIMESTAMPTZ`
- `LOCALTIMESTAMP` for a `TIMESTAMP`
  
- `CURRENT_DATE` for a `DATE`
- `LOCALTIME` for a `TIME`



# CHAR(n) / VARCHAR(n)

Padded with whitespace up to length n

- Padding spaces are ignored when comparing
  - But not for pattern matching with **LIKE** & regular expressions!
- Actually not stored as fixed-width field!
  - Can waste space storing irrelevant spaces
  - Performance-wise, spend extra time stripping spaces
  - Index created for **CHAR(n)** may not work with a **TEXT** parameter
- **company\_name VARCHAR(50)** → Peterson's and Sons and Friends Bits & Parts Limited
- To restrict length, just enforce **CHECK** constraint
- Bottom line: just use **TEXT (VARCHAR)**



# MONEY

## Get away

- Fixed-point
  - Doesn't handle fractions of a cent, etc. – rounding may be off!
- Doesn't store currency type, assumes server **LC\_MONETARY**

- Accepts garbage input:

```
# SELECT ',123,456,,7,8.1,0,9'::MONEY;  
money  
-----  
£12,345,678.11  
(1 row)
```

- Just use **NUMERIC** and store currency in another column





# SERIAL

Used to be useful shorthand but now more trouble than it's worth

- Non SQL Standard
- Permissions for sequence created by **SERIAL** need to be managed separately from the table
- **CREATE TABLE ... LIKE** will use the same sequence!
- Use identity columns instead:

```
CREATE TABLE tab (id BIGINT GENERATED BY DEFAULT AS IDENTITY PRIMARY KEY,  
content TEXT);
```

- **With an identity column, you don't need to know the name of the sequence:**  
**ALTER TABLE** tab **ALTER COLUMN** id **RESTART WITH** 1000;
- **BUT:** if application depends on a serial sequence with no gaps (e.g. for receipt numbers), generate that in the application



# Improper feature usage

# SQL\_ASCII

## Is not a database encoding

- No encoding conversion or validation!
  - Byte values **0-127** interpreted as ASCII
  - Byte values **128-255** uninterpreted
- Setting behaves differently from other character sets
- Can end up storing a mixture of encodings
  - And no way to recover original strings



# CREATE RULE

## RULEs are not the same as TRIGGERs

- Rules don't simply apply conditional logic
  - They rewrite queries to modify or add extra queries
  - All non-trivial rules will probably have unintended side-effects
  - Non SQL Standard
- If you are not creating writable **VIEWs**, use **TRIGGERs** instead
- Look for Depesz's exhaustive blog post on rules:  
<https://www.depesz.com/2010/06/15/to-rule-or-not-to-rule-that-is-the-question>



# CREATE TABLE (...) INHERITS ...

(i)

## Table inheritance

- Seemed like a good idea before ORMs...
- e.g. 

```
CREATE TABLE events (id BIGINT, ... many columns ... );  
CREATE TABLE meetings (scheduled_time TIMESTAMPTZ)  
INHERITS events;
```
- Was used to implement partitioning (< PG 10)
- Incompatible with declarative partitioning (>= PG 10):
  - One cannot inherit from a partitioned table
  - One cannot add inheritance to a partitioned table



# CREATE TABLE (...) INHERITS ...

(ii)

## How to undo table inheritance

- You can replace table inheritance with foreign key relations
- Create a new table to hold the data, and add the FK column:

```
CREATE TABLE new_meetings LIKE meetings;  
ALTER TABLE new_meetings ADD item_id BIGINT;
```

- Copy data from old table into new one (may take a long time):

```
INSERT INTO new_meetings  
SELECT *, id FROM meetings;
```

- Create required constraints, indexes, triggers etc. for new\_meetings



# CREATE TABLE (...) INHERITS ...

(iii)

## How to undo table inheritance (continued)

- **Very dirty hack** (if your table is huge) - create the FK but do not validate it now to avoid the full table scan:

```
ALTER TABLE new_meetings
CONSTRAINT event_id_fk
FOREIGN KEY (event_id)
REFERENCES events (id)
NOT VALID;
```

- If doing this on a live system, create a trigger to replicate changes coming into **meetings** also into **new\_meetings**
- Normally one should not touch **pg\_catalog** directly, but we can **UPDATE pg\_constraint SET convalidated = true WHERE conname = 'event\_id\_fk'**; as we are confident that data in FK column is valid (as exact copy of the original table)



# CREATE TABLE (...) INHERITS ...

(iv)

## How to undo table inheritance (continued)

- Inside a transaction, perform all the DDL at once:

```
DO $$  
BEGIN  
  ALTER TABLE meetings RENAME TO old_meetings;  
  ALTER TABLE new_meetings RENAME TO meetings;  
  DROP TABLE old_meetings;  
  -- IMPORTANT: Create trigger to INSERT/UPDATE/DELETE items in  
  -- events as they get changed in meetings - it's easy as now  
  -- we have the FK.  
  COMMIT;  
END $$ LANGUAGE plpgsql;
```





# Partitioning by multiple keys

(i)

## Is not partitioning on multiple levels

- Be careful!
- `CREATE TABLE transactions ( ... , location_code TEXT, tstamp TIMESTAMPTZ)  
PARTITION BY RANGE (tstamp, location_code);`
- `CREATE TABLE transactions_2023_02_a  
PARTITION OF transactions  
FOR VALUES FROM ('2023-02-01', 'AAA') TO ('2023-03-01', 'BAA');`
- `CREATE TABLE transactions_2023_02_b  
PARTITION OF transactions  
FOR VALUES FROM ('2023-02-01', 'BAA') TO ('2023-03-01', 'BZZ');`  
`ERROR: partition "transactions_2023_02_b" would overlap partition  
"transactions_2023_02_a"`



# Partitioning by multiple keys

(ii)

Subpartitioning is what you actually need

- `CREATE TABLE transactions ( ... , location_code TEXT, tstamp TIMESTAMPTZ)  
PARTITION BY RANGE (tstamp);`
- `CREATE TABLE transactions_2023_02  
PARTITION OF transactions  
FOR VALUES FROM ('2023-02-01') TO ('2023-03-01')  
PARTITION BY HASH (location_code);`
- `CREATE TABLE transactions_2023_02_p1  
PARTITION OF transactions_2023_02  
FOR VALUES WITH (MODULUS 4, REMAINDER 0);`



# Performance considerations

# Number of connections

(i)

Don't overload your server for no reason

- `max_connections = 5000`
- Every client connection spawns a separate backend process
  - IPC via semaphores & shared memory
  - Risk: CPU context switching
- Accessing the same objects from multiple connections may incur many Lightweight Locks (LWLocks or “latches”)
  - Lock becomes heavily contended, lots of lockers slow each other down
  - You may be making your data hotter for no reason
  - No queuing, more or less random



# Number of connections

(ii)

## Mitigation strategy

- Pre-PG 13: **Snapshot contention**
  - Each transaction has an MVCC snapshot – even if idle!
- Contention often caused by too much concurrency
  - Insert a connection pooler (e.g. **PgBouncer**) between application and DB
  - Allow fewer connections into the DB, make the rest queue for their turn
  - “Throttle” or introduce latency on the application side, to save your server performance
    - Sounds counter-intuitive!
    - Doesn’t necessarily slow anything down – queries may execute faster!



# High transaction rate

(i)

Just because you can, doesn't mean you should

- Postgres assigns an identifier to each transaction
  - Unsigned 32-bit int (4.2B values)
  - Circular space, with a visibility horizon
- **XID wraparound**: you try to read a very old tuple that is > 2.1B XIDs in the past
- Very heavy OLTP workloads can go through 2.1B transactions in a short time
  - For you, that's the future! (invisible)
  - **Freezing**: Flag tuple as “frozen” which is known to always be in the past
- Need to make sure **FREEZE** happens before XID wraparound



# High transaction rate

(ii)

## What can you do?

- Can batching help?
  - Does application really need to commit everything atomically?
  - Batch size 1000 will have 1/1000th the burn rate
- Increase effectiveness of autovacuum
  - More efficient FREEZE



# Turning off autovacuum

(i)

a.k.a. the MVCC maintenance operation. Yeah, don't.

- Removes dead tuples, freezes tuples (among other things)
- Has overhead
  - Scans tables & indexes
  - Needs, obtains, and waits for locks
  - Has limited capacity by default
- People are concerned about overhead
  - Alternative is worse! You can't avoid **VACUUM** in Postgres (yet).
  - You can outrun it (and then you'll need **VACUUM FULL**)





# Turning off autovacuum

(ii)

For most production workloads, defaults are too low

- Make it work harder to avoid problems
- Increase potency via:
  - `maintenance_work_mem` (1GB is good)
  - `autovacuum_max_workers`
  - `autovacuum_vacuum_cost_delay` / `autovacuum_vacuum_cost_limit`



# Explicit locking

(i)

## a.k.a. heavyweight locks

- Table-level (e.g. **SHARE**) or row-level (e.g. **FOR UPDATE**)
- Conflict with other lock modes (e.g. **ACCESS EXCLUSIVE** with **ROW EXCLUSIVE**)
- Block read/write access totally leading to waits
- Disastrous for performance
  - Unless your application is exquisitely crafted
  - Hint: it isn't



# Explicit locking

(ii)

## Lock contention: waiting for explicit locks

- Avoid explicit locking!
- Use SSI (Serializable Snapshot Isolation, **SERIALIZABLE** isolation level)
- Make application tolerant
  - Allow it to fail and retry
- Slightly reduced concurrency, but:
  - No blocking, no explicit locks needed (SIReadLocks, rw-conflicts)
  - Best performance choice for some application types



# Security considerations

# psql --W or --password

Request password before attempting connection

- It will ask for a password even if the server doesn't require one
- Unnecessary: `psql` will always ask for a password if required by server
- Insecure: You may think you're logging in with a password
  - But the server may be in `trust` mode and letting you in anyhow
  - Also, you may be entering the wrong password and still getting in
  - From a different client, you may get a surprise!



# listen\_addresses = "\*"

## Listening for connections from clients

- There's a **reason** the default is **'localhost'** (only TCP/IP loopback)
- Make sure you only enable the interfaces and networks which you actually want to have access to the database server
- e.g. Internet connection on one network & private network on another interface
- Don't advertise your presence:  
**3,600,000** MySQL/MariaDB servers (port 3306) found exposed on the Internet in May 20

## Come In WE'RE OPEN



# pg\_hba.conf → trust

## Host-Based Authentication

- Called that for a reason, i.e. configuring with `host ...` like:

```
host mydb myuser 10.10.10.10/32 md5
```

- `trust` with `host(ssl)` is a **Very Bad Idea**
  - Even for `local` e.g. improper user can connect to the DB
  - Postgres might be fine, but other software on the same server could be compromised
- Default to giving access only where strictly necessary (better safe...)

## No door



# Database owned by superuser

## Do you really need to?

- Use superuser only for management of global objects
  - Such as users
  - Good security practice
- Superuser bypasses a lot of checks
- (Bad) code that's normally harmless could be exploited in harmful way with superuser access
- Try to restrict database ownership to standard users





# Thank you!

Find me on Mastodon: [@vyruss@fosstodon.org](https://fosstodon.org/@vyruss)

Photo: “The Devil’s Beef Tub”, Scotland

# Thank you!



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